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Boolean inverse monoids

Mark V. Lawson

Heriot-Watt University

Coimbra, Portugal

June 2019

With the collaboration of many

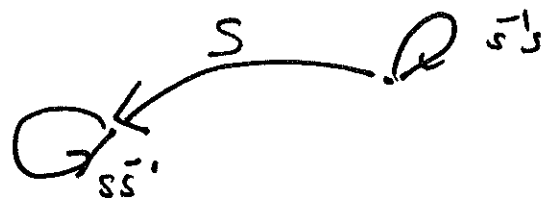
colleagues

- We deal here only with monoids which corresponds to working with compact spaces. Our results can also be extended to semigroups which corresponds to working with locally compact spaces
- The theory generalizes to distributive inverse semigroups and pseudogroups

Recollections of some terminology

Let S be an inverse semigroup

- If $s \in S$ write $d(s) = \bar{s}s$ and $r(s) = s\bar{s}$. Regard s as an arrow



- $E(S)$ is the semilattice of idempotents of S .
- \leq denotes the natural partial order defined by $a \leq b$ iff $a = be$ for some idempotent e .

- \sim denotes the compatibility relation defined by $a \sim b$ iff $\bar{a}'b$ and $a\bar{b}'$ are both idempotents.
- \perp denotes the orthogonality relation defined by $a \perp b$ iff $\bar{a}'b = 0 = a\bar{b}'$.
- An inverse monoid is factorizable if every element is beneath a unit (w.r.t. the natural partial order).
- An inverse semigroup is fundamental if the only elements that commute with all idempotents are idempotents.

Key definitions

The inverse resid S is said to be distributive iff:

- $a \sim b \Rightarrow \exists a \vee b.$

- If $\exists a \vee b$ then

$$c(a \vee b) = ca \vee cb \quad \text{and}$$

$$(a \vee b)c = ac \vee bc \quad (\forall c \in S)$$

When S is distributive, $E(S)$ is a distributive lattice.

An inverse monoid is said to be Boolean if it is distributive and its semilattice of idempotents is a Boolean algebra.

Morphisms of Boolean inverse monoids are monoid homomorphisms that preserve binary compatible joins and restrict to Boolean algebra maps between the respective Boolean algebras of idempotents.

Notation If $b \leq a$ define

$$a \setminus b = a \overline{d(b)}$$

Example Let $X \neq \emptyset$ be any set.

Then $I(X)$ denotes the symmetric inverse monoid on X . It is clear that $I(X)$ is a Boolean inverse monoid. We shall usually denote this monoid by $I|X|$.

Thus I_5 is the symmetric inverse monoid on the set $X = \{1, 2, 3, 4, 5\}$.

Direct products of Boolean inverse monoids are Boolean inverse monoids.

The $I_{n_1} \times \dots \times I_{n_r}$

is a Boolean inverse monoid.

These are Boolean inverse monoids which are analogues of AF, Cuntz, and Cuntz-Krieger C^* -algebras

Boolean inverse monoids arise naturally in ring theory (and C^* -algebras).

Theorem Let S be an inverse submonoid (with zero) of the multiplicative monoid of a ring R . Then there is a Boolean inverse submonoid S'' of R s.t.

$$S \subseteq S'' \subseteq R.$$

In fact, Boolean inverse monoids can be constructed from any inverse monoid.

Theorem (Booleanization) Let S

be any inverse monoid with zero.

Then there is a Boolean inverse monoid

$\underline{B}(S)$ together with an embedding

$$\beta: S \rightarrow \underline{B}(S)$$

s.t. if $\theta: S \rightarrow T$ is any homomorphism to a Boolean inverse semigroup then

there is a unique morphism $\gamma: \underline{B}(S) \rightarrow T$

$$\text{s.t. } \theta = \beta\gamma.$$

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Boolean inverse monoids are more
"ring-like" than arbitrary ~~inverse~~ inverse
monoids.

Let S be a Boolean inverse monoid.
 An additive ideal $I \triangleleft S$ is a
 semigroup ideal I closed under
 binary compatible joins.

Lemma

- (1) The kernels of morphisms between Boolean inverse monoids are additive ideals.
- (2) The kernel of a morphism is trivial
 \Leftrightarrow that morphism is idempotent-separating.

We can use ideals to build quotients
(but not Rees quotients).

Lemma Let $I \trianglelefteq S$ be an
additive ideal. Define \equiv_I on S
by $a \equiv_I b \Leftrightarrow (\exists c \leq a, b);$

$$(a \wedge c), (b \wedge c) \in I.$$

Then $S / \equiv_I = S / I$ is a

Boolean inverse monoid and the natural
map $S \rightarrow S / I$ is a morphism with
kernel I .

A congruence equal to some \equiv_I is
said to be ideal-induced.

N.B. Not every morphism is ideal-induced.

Anja Kudryartseva proved that

they are precisely the weakly meet

preserving ones. That if $\theta: S \rightarrow T$

then it is weakly meet preserving if

$t \leq \theta(a), \theta(b) \rightarrow \exists c \leq a, b$ s.t.

$t \leq \theta(c)$.

0-simplifying = no, non-trivial additive ideals

Simple = 0-simplifying + fundamental
(no additive congruences)

Strongly simple = 0-simple + fundamental
(no congruences of any description)

Program: describe all simple Boolean inverse monoids.

This program may seem overambitious, but there are clues.

An element a of a Boolean inverse monoid is finite if the number of elements below it is finite.

A Boolean inverse monoid is said to be semi simple if every element is finite.

An atom is an element a s.t.

$$b \leq a \Rightarrow b = a \text{ or } b = 0.$$

The dichotomy theorem Let S

be a 0 -simplifying Boolean inverse monoid. Then exactly one of the following holds:

- (1) S is atomless.
- (2) S is semisimple.

The class of countable, simple Boolean inverse monoids looks particularly interesting.

Three theorems



#1

Let $n \geq 1$ be a natural number.

Let G^0 be a group with zero adjoined.

$R_n(G^0)$ denotes the set of all $n \times n$ rook matrices over G^0 . A rook matrix is a square matrix with at most one non-zero entry in each row and each column.

Theorem 1

- (1) The finite Boolean inverse monoids have the form $R_{n_1}(G_1^0) \times \dots \times R_{n_r}(G_r^0)$.
- (2) The finite fundamental Boolean inverse monoids have the form $I_{n_1} \times \dots \times I_{n_r}$.
- (3) The finite simple Boolean inverse monoids have the form I_n .

#2

Let S be a Boolean inverse monoid.
 Observe that V is partially defined.

Put

$$\cdot x \odot y = \varepsilon(x) \overline{\varepsilon(y)} = x \wedge \overline{d(y)}$$

[$x \odot y$ is the largest element $\leq x$
 and \perp to y]. Called skew difference.

$$\cdot x \nabla y = (x \odot y) \vee y.$$

Called left skew join.

Both \odot and ∇ are globally defined.

Theorem 2 (Wehrung) Boolean

inverse monoids can be axiomatized
by means of equations using the
signature $(S, 1, 0, ^{-1}, \cdot, \emptyset, \nabla)$.

Thus they form a variety.

In addition, there is a Mal'cev

term for this variety and so

they are Congruence-permutable.

#3

An MV algebra $(A, \oplus, \neg, 0)$ is a set equipped with a binary operation \oplus , a unary operation \neg , and a constant 0 s.t.

(MV1) \oplus is associative.

(MV2) \oplus is commutative.

(MV3) 0 is the identity for \oplus .

(MV4) $\neg\neg x = x$.

(MV5) $x \oplus \neg 0 = \neg 0$. Define

$1 = \neg 0$.

(MV6) $\neg(\neg x \oplus y) \oplus y =$

$\neg(\neg y \oplus x) \oplus x$.

MV-algebras arise as Lindenbaum algebras of Many-valued logics in the same way that Boolean algebras arise as Lindenbaum algebras of Classical, two-valued logic.

The type monoid of a Boolean
inverse monoid.

Let S be a Boolean inverse monoid.

$E(S) / \mathcal{D}$ is the set of \mathcal{D} -classes
of idempotents. Denote the \mathcal{D} -class
containing e by $[e]$.

Define $[e] \oplus [f] = [e' \vee f']$
if $e' \in [e]$, $f' \in [f]$, $e' \perp f'$.
(undefined otherwise).

Put $\underline{0} = [0]$ and $\underline{1} = [1]$.

A Commutative monoid M is said to be a refinement monoid if $a_1 + a_2 = b_1 + b_2$ implies there exist elements $c_{11}, c_{12}, c_{21}, c_{22}$

$$\text{s.t. } a_1 = c_{11} + c_{12}, \quad a_2 = c_{21} + c_{22},$$

$$b_1 = c_{11} + c_{21} \quad \text{and} \quad b_2 = c_{12} + c_{22}.$$

	b_1	b_2
a_1	c_{11}	c_{12}
a_2	c_{21}	c_{22}

Theorem Let S be a Boolean inverse monoid. Then the partial structure $(E(S)/\mathcal{D}, \oplus)$ can be completed in a unique way to a commutative refinement monoid $\text{Typ}(S)$, called the type monoid of S .

Lemma A Boolean inverse monoid is factorizable $\Leftrightarrow (e \in D_f \Rightarrow \bar{e} \in D_{\bar{f}})$.

A Foulis monoid is a factorizable, Boolean inverse monoid in which S/\mathcal{I} is a lattice.

In a Foulis monoid, we may define $\neg [e] = [\bar{e}]$.

Theorem 3 (Lawson & Scott, Wehrung)

Every MV algebra is isomorphic to the type monoid of some Foulis monoid.

Non - commutative

Stone duality

Classical Stone duality

A Boolean space is a compact, Hausdorff space with a basis of clopen subsets.

Stone duality establishes a dual equivalence between the category of Boolean spaces and the category of Boolean algebras.

Let B be a Boolean algebra.

Put

$\underline{X}(B)$ = set of ultrafilters of X .

If $a \in B$, put

U_a = ultrafilters that contain a .

Put

$$\beta = \{U_a : a \in B\}.$$

The β is the basis of a topology on $\underline{X}(B)$

that makes it a Boolean space.

Let X be a Boolean space.

Put $\underline{B}(X) =$ set of clopen subsets of X . The $\underline{B}(X)$ is a

Boolean algebra.

Theorem

(1) If B is a Boolean algebra
 $\kappa \quad B \cong \underline{B}(X(B)).$

(2) If X is a Boolean space
 $\kappa \quad X \cong \underline{X}(\underline{B}(X)).$

$\underline{X}(B)$ is called κ Stone space of B .

Example If X is a finite Boolean space then $\mathcal{B}(X)$ is endowed with the discrete topology. The $\mathcal{B}(X) = P(X)$, the Boolean algebra of all subsets of X . This is a finite Boolean algebra and every finite Boolean algebra is isomorphic to one of this kind.

Example There is, up to isomorphism,
exactly one countable, atomless Boolean
algebra. I shall call it the Tarski algebra.

The Stone space of the Tarski algebra
is the Cantor space.

Non-commutative Stone duality

We shall replace Boolean algebras
by Boolean inverse monoids.

Boolean spaces will be replaced
by a suitable class of
topological groupoids.

A groupoid G is a small category in which every arrow is invertible.



$G_0 \subseteq G$ is the set of identities $\neq G$, meaning elements e s.t. if $\exists ge$ then $ge = g$ and $\exists eg$ then $eg = g$.

A subset $A \subseteq G$ is called a partial biset if $a, b \in A$ and $\bar{a}a = \bar{b}b$ (resp. $a\bar{a} = b\bar{b}$) $\Rightarrow a = b$.

A groupoid G is equipped with maps

$$\underline{d}: G \rightarrow G, \text{ given by}$$

$$\underline{d}(g) = \bar{s}'g, \quad \underline{t}: G \rightarrow G, \text{ given by}$$

$$\underline{t}(g) = g\bar{s}' \text{ and a map } M: G * G \rightarrow G$$

given by $(g, h) \mapsto gh$ where

$$G * G = \{ (g, h) \in G * G : \underline{d}(g) = \underline{t}(h) \}.$$

and a map $G \rightarrow G$ given by $g \mapsto \bar{g}'$.

If G is equipped with a topology s.t. these maps are continuous we say

G is a topological groupoid.

A ~~groupoid~~ topological groupoid G is said to be étale if d is a local homeomorphism.

The algebraic significance of étale groupoids is provided by the following result of Pedro Resende.

Theorem Let G be an étale groupoid. The $\Omega(G)$, the set of all open subsets of G , is a mooid under product of subsets.

An étale topological groupoid is said to be Boolean if its space of identities is a Boolean space.

Idea

Boolean algebras

replaced by \longrightarrow

Boolean inverse monoids

Boolean spaces

replaced by \longrightarrow

Boolean groupoids.

Let G be a Boolean groupoid.

Denote by $\underline{KB}(G)$ the set of all
 compatible partial bisections of G .

Proposition Under subset multiplication,

$\underline{KB}(G)$ is a Boolean inverse

groupoid.

Let S be a Boolean inverse monoid.
Denote by $\underline{G}(S)$ the set of all
ultrafilters on S .

It is not immediately obvious
but the set $\underline{G}(S)$ can be
endowed with the structure of a
groupoid.

If $s \in S$, denote by U_s the set of
all ultrafilters in S that contain s .

Put $\beta = \{U_a : a \in S\}$. Then β is the
basis for a topology on $\underline{G}(S)$.

Proposition $\underline{G}(S)$ is a Boolean
groupoid. We call it the Stone groupoid
of S .

Theorem

(1) If S is a Boolean inverse monoid then $S \cong \underline{KB}(G(S))$.

(2) If G is a Boolean groupoid then $G \cong \underline{G}(\underline{KB}(G))$.

Refinements

- $\underline{G}(S)$ is Hausdorff $\Leftrightarrow S$ has binary meets.
- $\underline{G}(S)$ is effective $\Leftrightarrow S$ is fundamental.
- $G_1 \sqcup G_2$ corresponds to $\underline{KB}(G_1) \times \underline{KB}(G_2)$.
- $G_1 \times G_2$ corresponds to $\underline{KB}(G_1) \otimes \underline{KB}(G_2)$.

Theorem Let S be an
 inverse monoid with zero. Then the
 Stone groupoid of $\underline{B}(S)$, the
 Booleanization of S , is Paterson's
universal groupoid of S .

Example If G is a finite Boolean groupoid then it is endowed with the discrete topology. Thus every finite Boolean inverse monoid is isomorphic to the Boolean inverse monoid of all partial bisections of a finite discrete groupoid.

This result can be refined to give the structure theory of finite Boolean inverse monoids described earlier.

Example Second-countable Boolean
 groupoids whose space of identities is the
 Cantor space correspond to countably
 infinite Boolean inverse monoids with a
 Tarski algebra of idempotents.

These play an important role in modern
 group theory, see work of Matui et al,
 via their groups of units.

A Tarski monoid is a countably infinite, atomless Boolean inverse Λ -monoid.

Theorem (Matui, Boolean inverse monoid version). Let S and T be simple Tarski monoids. Then t.f.a.e.

(1) The monoids S and T are isomorphic.

(2) Their groups of units are isomorphic.