

Non-commutative Stone duality

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1. Classical Stone duality

A *Boolean algebra* is an algebraic structure $(B, +, \cdot, \bar{}, 0, 1)$ of type $(2, 2, 1, 0, 0)$ satisfying the following axioms:

$$(B1) \quad (x + y) + z = x + (y + z).$$

$$(B2) \quad x + y = y + x.$$

$$(B3) \quad x + 0 = x.$$

$$(B4) \quad (x \cdot y) \cdot z = x \cdot (y \cdot z).$$

$$(B5) \quad x \cdot y = y \cdot x.$$

$$(B6) \quad x \cdot 1 = x.$$

$$(B7) \quad x \cdot (y + z) = x \cdot y + x \cdot z.$$

$$(B8) \quad x + (y \cdot z) = (x + y) \cdot (x + z).$$

$$(B9) \quad x + \bar{x} = 1.$$

$$(B10) \quad x \cdot \bar{x} = 0.$$

Although this is an algebraic structure, we can also define a partial order by

$$y \leq x \Leftrightarrow y = y \cdot x.$$

An element x is called an *atom* if $y \leq x$ implies either that $y = x$ or $y = 0$.

A Boolean algebra is said to be *atomless* if it has no atoms.

A topological space X is called a *Boolean space* if it is compact, Hausdorff and 0-dimensional (that is, it has a base of clopen sets).

- If B is a Boolean algebra, then one can construct a Boolean space, $X(B)$, called the *Stone space* of B .
- If X is a Boolean space, then one can construct a Boolean algebra, $B(X)$, of the clopen subsets of X .

Theorem (Stone, 1937) Classical Stone duality.

1. *For any Boolean algebra B , we have that $B \cong \mathcal{B}(X(B))$.*
2. *For any Boolean space X , we have that $X \cong X(\mathcal{B}(X))$.*

Examples

1. If B is a finite Boolean algebra, then $X(B)$ is a finite set regarded as a discrete topological space consisting of all the atoms of B .
2. If B is a (=the) countably infinite atomless Boolean algebra (which I call the *Tarski algebra*) then $X(B)$ is the Cantor space.

Proofs There are two kinds of proof:

1. Use prime filters (=ultrafilters). See Givant and Halmos.
2. Use locales and then specialize. See Johnstone.

2. Ideas behind non-commutative Stone duality

1. Replace the Boolean algebra by some kind of semigroup which has a 'Boolean character'. The algebraic structure will be used to define a partial order. The \cdot of the Boolean algebra will become the semigroup operation, and the $+$ will become a partially defined join. The 'Boolean character' will be reflected in the fact that a set of idempotents will be required to form a Boolean algebra.
2. Replace the topological space by a (1-sorted) topological category. This is a small category equipped with a topology with respect to which the operations \mathbf{d} (dom), \mathbf{r} (cod), \mathbf{m} (multiplication) are required to be continuous. In addition, we require that the space of identities is open and forms a Boolean space.

3. Non-commutative Stone duality

The generalization of Boolean algebras.

We shall replace Boolean algebras by a class of monoids called Boolean right restriction monoids.

We shall motivate this abstract class of monoids with a concrete example.

Denote by $PT(X)$ the set of all partial functions defined on the (non-empty) set X .

An element of $PT(X)$ has the form $f: A \rightarrow X$ where $A \subseteq X$. We call the subset A the *domain of definition* of f denoted by $\text{dom}(f)$.

Denote by f^* the identity function defined on $\text{dom}(f)$. Whereas identity functions defined on subsets of X are idempotents, it is not true that all idempotents have this form. Idempotents that are identity functions defined on subsets are called *projections*. The set of projections of $PT(X)$ is denoted by $\text{Proj}(PT(X))$.

Partial functions f and g can be compared using subset inclusion. We say that $f \subseteq g$ precisely when g restricted to the domain of definition of f is f itself. In fact, $f \subseteq g$ if and only if $f = gf^*$. Thus we do not have to impose a partial order from the outside since the order can be defined purely algebraically.

If $f, g \in \text{PT}(X)$, it is not true in general that $f \cup g \in \text{PT}(X)$ since $f \cup g$ might not be a partial function. It is a partial function precisely when for all $x \in \text{dom}(f) \cap \text{dom}(g)$ we have that $f(x) = g(x)$. This means precisely that $f(f^*g^*) = g(f^*g^*)$ which simplifies to $fg^* = gf^*$. We shall say that f and g are *left-compatible* if $fg^* = gf^*$. The partial functions f and g are left-compatible precisely when $f \cup g$ is also a partial function.

Observe that the set of projections is a commutative, idempotent subsemigroup. In fact, the set of projections forms a Boolean algebra since it is isomorphic to the powerset of X with respect to subset inclusion.

There is a map $f \mapsto f^*$ from $\text{PT}(X)$ to $\text{Proj}(\text{PT}(X))$. Thus we regard $\text{PT}(X)$ as an algebra of type $(2, 1)$. Observe that $ff^* = f$ and that $f^*g = g(fg)^*$.

We now abstract the above example.

We define a monoid S to be a *right restriction monoid* if it is equipped with a unary operation $a \mapsto a^*$ satisfying the following axioms:

$$(RR1) \quad (s^*)^* = s^*.$$

$$(RR2) \quad (s^*t^*)^* = s^*t^*.$$

$$(RR3) \quad s^*t^* = t^*s^*.$$

$$(RR4) \quad ss^* = s.$$

$$(RR5) \quad (st)^* = (s^*t)^*.$$

$$(RR6) \quad t^*s = s(ts)^*.$$

The unary operation $s \mapsto s^*$ is called *star*. Denote by $\text{Proj}(S)$ those elements a such that $a^* = a$ and call them *projections*.

We define a left restriction monoid, dually.

A *bi-restriction monoid* is a monoid which is both a left and right restriction monoid and the sets of projections are the same.

Let S be a right restriction monoid. Define

$$y \leq x \text{ iff } y = xy^*.$$

This is a partial order with respect to which the monoid is partially ordered. This is called the *natural partial order*.

If $a, b \leq c$ then $ab^* = ba^*$. More generally, we say that a and b are *left compatible* if $ab^* = ba^*$. Being left compatible is therefore a necessary condition to have a join.

A *Boolean right restriction monoid* is a right restriction monoid that satisfies the following three conditions:

1. Left compatible elements have joins.
2. Multiplication distributes over any joins that exist.
3. The set of projections forms a Boolean algebra with respect to the natural partial order.

The monoids $\text{PT}(X)$ are Boolean right restriction monoids.

The generalization of Boolean spaces.

A topological category is said to be *domain-étale* if \mathbf{d} is a local homeomorphism. Dually, we talk about *range-étale*, and *étale* if both.

A domain-étale category is said to be *Boolean* if its space of identities is a Boolean space.

We now have all the ingredients we need for non-commutative Stone duality:

- Boolean restriction monoids.
- Boolean domain-étale topological categories.

We sketch out how to construct a Boolean domain-étale topological category $\mathcal{C}(S)$ from a Boolean restriction monoid S .

Put $\mathcal{C}(S)$ equal to the prime filters in S .

If A is a prime filter, put

$$\mathbf{d}(A) = (A^*)^\uparrow \text{ and } \mathbf{r}(A) = \{e \in \text{Proj}(S) : eA \subseteq A\}^\uparrow.$$

Both of these are also prime filters.

If A and B are prime filters and $\mathbf{d}(A) = \mathbf{r}(B)$ define

$$A \cdot B = (AB)^\uparrow.$$

This gives us a category on $\mathcal{C}(S)$.

If $a \in S$ let X_a be all the prime filters of S that contain a .

Put $\beta = \{X_a : a \in S\}$. This is the base for a topology on $C(S)$.

Theorem *If S is a Boolean right restriction monoid then $C(S)$ is a Boolean domain-étale topological category.*

We now go in the opposite direction.

Let C be a Boolean domain-étale topological category. A subset $A \subseteq C$ is called a *local section* if $a, b \in A$ and $d(a) = d(b)$ then $a = b$.

Denote by $KS(C)$ the set of all compact-open local sections of C .

Theorem *If C is a Boolean domain-étale topological category then $KS(C)$ is a Boolean right restriction monoid.*

Theorem Non-commutative Stone duality.

1. *If S is a Boolean right restriction monoid then $S \cong \text{KS}(C(S))$.*
2. *If C is a Boolean domain-étale topological category then $C \cong C(\text{KS}(C))$.*

This was first proved by Cockett and Garner (2021) using locales; the proof I have sketched using prime filters is Lawson (2024).

4. The étale case

Let S be a Boolean right restriction monoid. We can ask the question when $C(S)$ is a groupoid.

We say that $a \in S$ is a *partial unit* if there exists $b \in S$ such that $ab = b^*$ and $ba = a^*$.

The set of partial units is denoted by $\text{Inv}(S)$ and is a Boolean inverse monoid.

We say that a Boolean right restriction monoid is *étale* if each element is a finite join of left-compatible partial units.

Theorem (Garner) *Let S be a Boolean right restriction monoid. Then $C(S)$ is a groupoid iff S is étale.*

Theorem (Lawson, 2024) *The structure of an étale Boolean right restriction monoid S is completely determined by $\text{Inv}(S)$.*

5. Other special cases

- Boolean inverse monoids and Boolean étale topological groupoids are important in the theory of C^* -algebras. Resende, 2007, Lawson, 2010, and Lawson & Lenz, 2013. Proved using prime filters.
- Boolean bi-restriction monoids and Boolean étale topological categories. Kudryavtseva & Lawson, 2017. Proved using locales.

6. Applications

- The duality between Boolean inverse monoids and Boolean groupoids is bound up with the theory of a class of groups, called *topological full groups*, that includes the Thompson-Higman groups. See the papers by Lawson, Sims & Vdovina, 2020, 2024.
- *Garner proved that the category of Cartesian closed universal algebras is equivalent to the category of Boolean right restriction monoids.*

7. References

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